## **IN THE CLAIMS**:

Kindly amend the claims, as follows:

1. (Original) A data transmission system for transmitting user data to and receiving data from a communication channel, comprising:

a parity check matrix comprising:

M tiers,

wherein  $M \ge 2$ ,

Dmin = 2 \* M for M = 1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

tc = M

wherein tc is the column weight, and

cycle - 4 = 0;

a linear block encoder to encode the user data in response to said parity check matrix;

a transmitter to transmit an output of said linear block encoder to the communication channel;

a soft channel decoder to decode data; and

a soft linear block code decoder to decode data decoded by said soft channel decoder in response to said parity check matrix.

- 2. (Original) A data transmission system according to Claim 1, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 3. (Original) A data transmission system according to Claim 2, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
  - 4. (Original) A data transmission system according to Claim 2, wherein said M

tiers are arranged in increasing rank order.

- 5. (Original) A data transmission system according to Claim 4, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
- 6. (Original) A data transmission system according to Claim 2, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .
- 7. (Original) A data transmission system according to Claim 6, wherein said matrix comprises  $\sum_{i=1}^{M} P_i (M-1)$  independent rows.
- 8. (Original) A data transmission system according to Claim 2, wherein said matrix comprises  $\sum_{i=1}^{M} P_i (M-1)$  parity bits.
- 9. (Original) A data transmission system according to Claim 2, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- 10. (Original) A data transmission system according to Claim 4, wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, & \text{if } c \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, & \text{otherwise} \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 11. (Currently Amended) A data transmission system according to Claim 4, wherein M=3, the number of rows =  $P_1 + P_2 + P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and tc = 3.
- 12. (Currently Amended) A data transmission system according to Claim 11, wherein a code rate =  $(P_1 * P_2 + P_1 P_2 P_3 + 2)/(P_1 * P_2)$ .
- 13. (Original) A data transmission system according to Claim 1, wherein said linear block code encoder comprises a low-density parity-check encoder and wherein said soft linear block code decoder comprises a low-density parity-check decoder.
- 14. (Original) A data transmission system according to Claim 1, wherein said soft channel decoder comprises a soft Viterbi algorithm decoder.
- 15. (Original) A decoder for decoding data from a communication channel, comprising:

a parity check matrix comprising:

M tiers.

wherein  $M \ge 2$ ,

Dmin = 2 \* M for M = 1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

tc = M,

wherein tc is the column weight, and

cycle -4 = 0;

a soft channel decoder to decode data; and

a soft linear block code decoder to decode data decoded by said soft channel decoder in accordance with said parity check matrix.

16. (Original) A decoder according to Claim 15, wherein each of said M tiers

comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .

- 17. (Original) A decoder according to Claim 16, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 18. (Original) A decoder according to Claim 16, wherein said M tiers are arranged in increasing rank order.
- 19. (Original) A decoder according to Claim 18, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
- 20. (Original) A decoder according to Claim 16, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .
- 21. (Original) A decoder according to Claim 20, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 22. (Original) A decoder according to Claim 16, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 23. (Original) A decoder according to Claim 16, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.

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24. (Original) A decoder according to Claim 18, wherein for each element A<sub>r,c</sub>

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, & \text{if } c \operatorname{mod}(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, & \text{otherwise} \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 25. (Currently Amended) A decoder according to Claim 18, wherein M=3, the number of rows =  $P_1+P_2+P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and tc = 3.
- 26. (Currently Amended) A decoder according to Claim 25, wherein a code rate =  $(P_1 * P_2 + P_1 P_2 P_3 + 2)/(P_1 * P_2)$ .
- 27. (Original) A decoder according to Claim 15, wherein said linear block code encoder comprises a low-density parity-check encoder and wherein said soft linear block code decoder comprises a low-density parity-check decoder.
- 26 28. (Currently Amended) A decoder according to Claim 15, wherein said soft channel decoder comprises a soft Viterbi algorithm decoder.
- 27 29. (Currently Amended) An encoder for encoding data from a communication channel, comprising:

a parity check matrix comprising:

wherein 
$$M \ge 2$$
,

Dmin = 
$$2 * M$$
 for  $M = 1..3$  or

$$2*M \ge Dmin \ge 6$$
 for  $M > 3$ ,

wherein Dmin is the minimum Hamming distance,

$$tc = M$$
.

wherein tc is the column weight, and

cycle 
$$-4 = 0$$
;

- a linear block encoder to encode the user data in response said parity check matrix; and
- a transmitter to transmit an output of said linear block encoder to the communication channel.
- 28 30. (Currently Amended) An encoder according to Claim 27 29, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 29 31. (Currently Amended) An encoder according to Claim 28 30, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 30 32. (Currently Amended) An encoder according to Claim 28 30, wherein said M tiers are arranged in increasing rank order.
- 31 33. (Currently Amended) An encoder according to Claim 30 32, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
- $\frac{32}{34}$ . (Currently Amended) An encoder according to Claim  $\frac{28}{30}$ , wherein said matrix comprises R rows, wherein R =  $\sum_{i=1}^{M} P_i$ .
- 33 35. (Currently Amended) An encoder according to Claim 32 34, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
  - 34 36. (Currently Amended) An encoder according to Claim 28 30, wherein said

matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.

- 35 37. (Currently Amended) An encoder according to Claim 28 30, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- $36 \ \underline{38}$ . (Currently Amended) An encoder according to Claim  $30 \ \underline{32}$ , wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, & \text{if } c \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, & \text{otherwise} \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 37 39. (Currently Amended) An encoder according to Claim 30 32, wherein M = 3, the number of rows =  $P_1+P_2+P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and to = 3.
- 38 <u>40</u>. (Currently Amended) An encoder according to Claim 37 <u>39</u>, wherein a code rate =  $(P_1 * P_2 + P_1 P_2 P_3 + 2)/(P_1 * P_2)$ .
- 39 41. (Currently Amended) An encoder according to Claim 27 29, wherein said linear block code encoder comprises a low-density parity-check encoder and wherein said soft linear block code encoder comprises a low-density parity-check encoder.
- 40 42. (Currently Amended) An encoder according to Claim 27 29, wherein said soft channel encoder comprises a soft Viterbi algorithm encoder.

41 43. (Currently Amended) A parity check matrix <u>embodied in a medium</u> for one of a low-density parity-check encoder and a low-density parity-check decoder, comprising:

M tiers,

wherein  $M \ge 2$ ,

Dmin = 2 \* M for M = 1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

tc = M,

wherein tc is the column weight, and

cycle -4 = 0.

- 42 <u>44</u>. (Currently Amended) A parity check matrix according to Claim 41 <u>43</u>, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 43 45. (Currently Amended) A parity check matrix according to Claim 42 44, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- [[44]] <u>46</u>. (Currently Amended) A parity check matrix according to Claim 42 <u>44</u>, wherein said M tiers are arranged in increasing rank order.
- 45 <u>47</u>. (Currently Amended) A parity check matrix according to Claim [[44]] <u>46</u>, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
- 46 <u>48</u>. (Currently Amended) A parity check matrix according to Claim 42 <u>44</u>, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .
  - 47 49. (Currently Amended) A parity check matrix according to Claim 46 48,

wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.

- 48 50. (Currently Amended) A parity check matrix according to Claim 42 44, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 49 51. (Currently Amended) A parity check matrix according to Claim 42 44, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- $50 \underline{52}$ . (Currently Amended) A parity check matrix according to Claim [[44]]  $\underline{46}$ , wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, ifc \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, otherwise \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 51 53. (Currently Amended) A parity check matrix according to Claim [[44]] 46, wherein M = 3, the number of rows =  $P_1 + P_2 + P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and tc = 3.
- 52 54. (Currently Amended) A parity check matrix according to Claim 51 53, wherein a code rate =  $(P_1 + P_2 + P_1 P_2 P_3 + 2)/(P_1 + P_2)$ .
- 53 55. (Currently Amended) A data transmission system for transmitting user data to and receiving data from a communication channel, comprising:

a parity check matrix comprising:

M tiers,

wherein  $M \ge 2$ ,

Dmin = 2 \* M for M = 1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

tc = M,

wherein tc is the column weight, and

cycle -4 = 0;

linear block encoding means for encoding the user data in response to said parity check matrix;

transmitting means for transmitting an output of said linear block encoding means to the communication channel;

soft channel decoding means for decoding data; and

soft linear block code decoding means for decoding data decoded by said soft channel decoding means in response to said parity check matrix.

- 54 <u>56</u>. (Currently Amended) A data transmission system according to Claim <u>53 55</u>, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 55 57. (Currently Amended) A data transmission system according to Claim 54 56, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 56 58. (Currently Amended) A data transmission system according to Claim 54 56, wherein said M tiers are arranged in increasing rank order.
- 57 <u>59</u>. (Currently Amended) A data transmission system according to Claim <u>56 58</u>, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .

- 58 <u>60</u>. (Currently Amended) A data transmission system according to Claim 54 <u>56</u>, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .
- 59 61. (Currently Amended) A data transmission system according to Claim 58 60, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 60 62. (Currently Amended) A data transmission system according to Claim 54 56, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 61 63. (Currently Amended) A data transmission system according to Claim 54 56, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- 62 64. (Currently Amended) A data transmission system according to Claim 56 58, wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, ifc \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, otherwise \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 63 65. (Currently Amended) A data transmission system according to Claim 56 58, wherein M=3, the number of rows =  $P_1+P_2+P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and tc = 3.
  - 64 66. (Currently Amended) A data transmission system according to Claim 63 65,

wherein a code rate =  $(P_1 - P_2 + P_1 - P_2 - P_3 + 2)/(P_1 - P_2)$ .

- 65 67. (Currently Amended) A data transmission system according to Claim 53 55, wherein said linear block code encoding means comprises a low-density parity-check encoding means and wherein said soft linear block code decoding means comprises a low-density parity-check decoding means.
- 66 68. (Currently Amended) A data transmission system according to Claim 53 55, wherein said soft channel decoding means comprises a soft Viterbi algorithm decoding means.
- 67 69. (Currently Amended) A decoder for decoding data from a communication channel, comprising:

a parity check matrix comprising:

M tiers,

wherein  $M \ge 2$ ,

Dmin = 2 \* M for M = 1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

tc = M,

wherein tc is the column weight, and

cycle -4 = 0;

soft channel decoding means for decoding data; and

soft linear block code decoding means for decoding data decoded by said soft channel decoding means in accordance with said parity check matrix.

- 68 70. (Currently Amended) A decoder according to Claim 67 69, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
  - 69 71. (Currently Amended) A decoder according to Claim 68 70, wherein the rank

of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.

- 70 72. (Currently Amended) A decoder according to Claim 68 70, wherein said M tiers are arranged in increasing rank order.
- 71 73. (Currently Amended) A decoder according to Claim 70 72, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
- 72 74. (Currently Amended) A decoder according to Claim 68 70, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .
- 73 75. (Currently Amended) A decoder according to Claim 72 74, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 74 76. (Currently Amended) A decoder according to Claim 68 70, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 75 77. (Currently Amended) A decoder according to Claim 68 70, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.

76  $\overline{78}$ . (Currently Amended) A decoder according to Claim 70  $\overline{72}$ , wherein for each element  $A_{r,c}$ 

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, & \text{if } c \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, & \text{otherwise} \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

77 79. (Currently Amended) A decoder according to Claim 70 72, wherein M = 3, the number of rows =  $P_1+P_2+P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and to = 3.

78 80. (Currently Amended) A decoder according to Claim 77 79, wherein a code rate =  $(P_1*P_2+P_1-P_2-P_3+2)/(P_1*P_2)$ .

- 79 81. (Currently Amended) A decoder according to Claim 67 69, wherein said linear block code encoding means comprises a low-density parity-check encoding means and wherein said soft linear block code decoding means comprises a low-density parity-check decoding means.
- 80 82. (Currently Amended) A decoder according to Claim 67 69, wherein said soft channel decoding means comprises a soft Viterbi algorithm decoding means.
- 81 83. (Currently Amended) An encoder for encoding data from a communication channel, comprising:

a parity check matrix comprising:

M tiers,

wherein  $M \ge 2$ ,

Dmin = 2 \* M for M = 1..3 or

$$2*M \ge Dmin \ge 6$$
 for  $M > 3$ ,

wherein Dmin is the minimum Hamming distance,

$$tc = M$$

wherein to is the column weight, and

cycle 
$$-4 = 0$$
;

linear block encoding means for encoding the user data in response said parity check matrix; and

transmitting means for transmitting an output of said linear block encoding means to the communication channel.

- 82 84. (Currently Amended) An encoder according to Claim 81 83, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 83 85. (Currently Amended) An encoder according to Claim 82 84, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 84 <u>86</u>. (Currently Amended) An encoder according to Claim <u>82</u> <u>84</u>, wherein said M tiers are arranged in increasing rank order.
- 85 87. (Currently Amended) An encoder according to Claim 84 86, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
- 86 88. (Currently Amended) An encoder according to Claim 82 84, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .
- 87 89. (Currently Amended) An encoder according to Claim 86 88, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.

- 88 90. (Currently Amended) An encoder according to Claim 82 84, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 89 91. (Currently Amended) An encoder according to Claim 82 84, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- 90 92. (Currently Amended) An encoder according to Claim 84 86, wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, & \text{if } c \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, & \text{otherwise} \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 91 93. (Currently Amended) An encoder according to Claim 84 86, wherein M = 3, the number of rows =  $P_1+P_2+P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and to = 3.
- 92 94. (Currently Amended) An encoder according to Claim 91 93, wherein a code rate =  $(P_1 * P_2 + P_1 P_2 P_3 + 2)/(P_1 * P_2)$ .
- 93 95. (Currently Amended) An encoder according to Claim 81 83, wherein said linear block code encoding means comprises a low-density parity-check encoding means and wherein said soft linear block code encoding means comprises a low-density parity-check encoding means.

- 94 96. (Currently Amended) An encoder according to Claim 81 83, wherein said soft channel encoding means comprises a soft Viterbi algorithm encoding means.
- 95 97. (Currently Amended) A method for transmitting data to and receiving data from a communication channel, comprising the steps of:
  - (a) generating a parity check matrix comprising:

M tiers,

wherein  $M \ge 2$ ,

Dmin = 2 \* M for M = 1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

tc = M,

wherein tc is the column weight, and

cycle - 4 = 0;

- (b) linear block encoding the data in accordance with the parity check matrix generated in step (a);
  - (c) transmitting the data encoded in step (b) to the communication channel;
  - (d) receiving the data from [[to]] the communication channel;
- (e) soft channel decoding the data read in step (d) in accordance with data decoded in step (g);
- (f) generating an address in accordance with the data soft linear block code decoding the data decoded in step (e); and
- (g) soft linear block code decoding data decoded by in step (e) in accordance with the address generated in step(f).
- 96 98. (Currently Amended) A method according to Claim 95 97, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 97 99. (Currently Amended) A method according to Claim 96 98, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said

identity matrix of another one of said tiers.

- 98 100. (Currently Amended) A method according to Claim 96 98, wherein said M tiers are arranged in increasing rank order.
- 99 101. (Currently Amended) A method according to Claim 98 100, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
- 100 102. (Currently Amended) A method according to Claim 96 98, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .
- 101 103. (Currently Amended) A method according to Claim 100 102, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 102 104. (Currently Amended) A method according to Claim 96 98, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 103 105. (Currently Amended) A method according to Claim 96 98, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.

 $\frac{104}{106}$ . (Currently Amended) A method according to Claim 98  $\frac{100}{100}$ , wherein for each element  $A_{r,c}$ 

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, & \text{if } c \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, & \text{otherwise} \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

 $\frac{105}{107}$ . (Currently Amended) A method according to Claim 98  $\frac{100}{100}$ , wherein M = 3, the number of rows = P<sub>1</sub>+P<sub>2</sub>+P<sub>3</sub>, the number of columns = P1 \* P2,  $\frac{100}{100}$  Pmin = 6, and tc = 3.

 $\frac{106}{108}$ . (Currently Amended) A method according to Claim  $\frac{105}{107}$ , wherein a code rate =  $(P_1*P_2+P_1-P_2-P_3+2)/(P_1*P_2)$ .

107 109. (Currently Amended) A method for decoding data received from a communication channel, comprising the steps of:

- (a) soft channel decoding the data received <u>from the communication channel</u> in accordance with data decoded in step (c);
  - (b) generating a parity check matrix comprising:M tiers,

wherein 
$$M \ge 2$$
,

Dmin = 
$$2 * M$$
 for  $M = 1..3$  or

$$2*M \ge Dmin \ge 6$$
 for  $M > 3$ ,

wherein Dmin is the minimum Hamming distance,

$$tc = M$$

wherein tc is the column weight, and

cycle 
$$-4 = 0$$
; and

(c) soft linear block code decoding data decoded [[by]] in step (a) in accordance

with the matrix generated in step(b).

- 108 110. (Currently Amended) A method according to Claim 107 109, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 109 111. (Currently Amended) A method according to Claim 108 110, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 110 112. (Currently Amended) A method according to Claim 108 110, wherein said M tiers are arranged in increasing rank order.
- 111 113. (Currently Amended) A method according to Claim 110 112, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
- 112 114. (Currently Amended) A method according to Claim 108 110, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .
- 113 115. (Currently Amended) A method according to Claim 108 114, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 114 116. (Currently Amended) A method according to Claim 108 110, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
  - 115 117. (Currently Amended) A method according to Claim 108 110, wherein

said matrix comprises  $P_1 \times P_2 - \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.

 $\frac{116}{118}$ . (Currently Amended) A method according to Claim  $\frac{110}{112}$ , wherein for each element  $A_{r,c}$ 

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, ifc \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, otherwise \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 117 119. (Currently Amended) A method according to Claim 110 112, wherein M = 3, the number of rows =  $P_1 + P_2 + P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and to = 3.
- 118 120. (Currently Amended) A method according to Claim 117 119, wherein a code rate =  $(P_1 P_2 + P_1 P_2 P_3 + 2)/(P_1 P_2)$ .
- 119 121. (Currently Amended) A method for encoding data transmitted to a communication channel, comprising the steps of:
  - (a) generating a parity check matrix comprising:M tiers,

wherein 
$$M \ge 2$$
,

Dmin = 
$$2 * M$$
 for  $M = 1..3$  or

$$2*M \ge Dmin \ge 6$$
 for  $M > 3$ ,

wherein Dmin is the minimum Hamming distance,

$$tc = M$$

wherein tc is the column weight, and

cycle 
$$-4 = 0;[[;]]$$

- (b) linear block encoding the data in accordance with the matrix generated in step (a); and
  - (c) transmitting the data encoded in step (b) to the communication channel.
- $\frac{120}{122}$ . (Currently Amended) A method according to Claim  $\frac{119}{121}$ , wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 121 123. (Currently Amended) A method according to Claim 120 122, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 122 124. (Currently Amended) A method according to Claim 120 122, wherein said M tiers are arranged in increasing rank order.
- $\frac{123}{25}$ . (Currently Amended) A method according to Claim  $\frac{122}{24}$ , wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
- 124 126. (Currently Amended) A method according to Claim 120 122, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .
- 125 127. (Currently Amended) A method according to Claim 124 126, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 126 128. (Currently Amended) A method according to Claim 120 122, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.

- 127 129. (Currently Amended) A method according to Claim 120 122, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- $\frac{128\ 130}{130}.$  (Currently Amended) A method according to Claim  $\frac{122\ 124}{124},$  wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, ifc \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, otherwise \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 129 131. (Currently Amended) A method according to Claim 122 124, wherein M = 3, the number of rows =  $P_1 + P_2 + P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and to = 3.
- $\frac{130 \text{ } 132}{132}$ . (Currently Amended) A method according to Claim  $\frac{129 \text{ } 131}{132}$ , wherein a code rate =  $(P_1 + P_2 + P_1 P_2 P_3 + 2)/(P_1 + P_2)$ .
- 131 133. (Currently Amended) A computer program embodied in a medium for transmitting data to and receiving data from a communication channel, comprising the steps of:
  - (a) generating a parity check matrix comprising:

M tiers,

wherein  $M \ge 2$ ,

Dmin = 2 \* M for M = 1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

tc = M

wherein tc is the column weight, and cycle -4 = 0;

- (b) linear block encoding the data in accordance with the parity check matrix generated in step (a);
  - (c) transmitting the data encoded in step (b) to the communication channel;
  - (d) receiving the data from [[to]] the communication channel;
- (e) soft channel decoding the data read in step (d) in accordance with data decoded in step (g);
- (f) generating an address in accordance with the data soft linear block code decoding the data decoded in step (e); and
- ( $\frac{1}{2}$ g) soft linear block code decoding data decoded by in step (e) in accordance with the address generated in step(f).
- 132 134. (Currently Amended) A computer program according to Claim 131 133, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 133 135. (Currently Amended) A computer program according to Claim 132 134, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 134 136. (Currently Amended) A computer program according to Claim 132 134, wherein said M tiers are arranged in increasing rank order.
- 135 137. (Currently Amended) A computer program according to Claim 134 136, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
- 136 138. (Currently Amended) A computer program according to Claim 132 134, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .

- 137 139. (Currently Amended) A computer program according to Claim 136 138, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 138 140. (Currently Amended) A computer program according to Claim 132 134, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 139 141. (Currently Amended) A computer program according to Claim 132 134, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- $\frac{140 \text{ } 142}{136}$ . (Currently Amended) A computer program according to Claim  $\frac{134}{136}$ , wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, ifc \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, otherwise \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 141 143. (Currently Amended) A computer program according to Claim 134 136, wherein M = 3, the number of rows =  $P_1 + P_2 + P_3$ , the number of columns=  $P_1 * P_2$ , dmin  $P_1 * P_2 * P_3$ , and to  $P_1 * P_2 * P_3$ , the number of columns=  $P_1 * P_2 * P_3$ , dmin  $P_1 * P_2 * P_3 * P_3 * P_4 * P_4 * P_5 * P_5 * P_6 * P_$
- 142 144. (Currently Amended) A computer program according to Claim 131 143, wherein a code rate =  $(P_1 * P_2 + P_1 P_2 P_3 + 2)/(P_1 * P_2)$ .
  - 143 145. (Currently Amended) A computer program embodied in a medium for

decoding data received from a communication channel, comprising the steps of:

- (a) soft channel decoding the data received <u>from the communication channel</u> in accordance with data decoded in step (c);
  - (b) generating a parity check matrix comprising:

M tiers,

wherein  $M \ge 2$ ,

Dmin = 2 \* M for M = 1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

tc = M,

wherein tc is the column weight, and

cycle -4 = 0; and

- (c) soft linear block code decoding data decoded [[by]] in step (a) in accordance with the matrix generated in step(b).
- 144 146. (Currently Amended) A computer program according to Claim 143 145, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 145 147. (Currently Amended) A computer program according to Claim 144 146, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 146 148. (Currently Amended) A computer program according to Claim 144 146, wherein said M tiers are arranged in increasing rank order.
- 147 149. (Currently Amended) A computer program according to Claim 146 148, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
  - 148 150. (Currently Amended) A computer program according to Claim 144

146, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .

- 149 151. (Currently Amended) A computer program according to Claim 148 150, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 150 152. (Currently Amended) A computer program according to Claim 144 146, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 151 153. (Currently Amended) A computer program according to Claim 144 146, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- $\frac{152}{154}$ . (Currently Amended) A computer program according to Claim  $\frac{146}{148}$ , wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, & \text{if } c \operatorname{mod}(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, & \text{otherwise} \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 153 155. (Currently Amended) A computer program according to Claim 146 148, wherein M = 3, the number of rows =  $P_1 + P_2 + P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and tc = 3.
- 154 156. (Currently Amended) A computer program according to Claim 153 155, wherein a code rate =  $(P_1 * P_2 + P_1 P_2 P_3 + 2)/(P_1 * P_2)$ .

- 155 157. (Currently Amended) A computer program embodied in a medium for encoding data transmitted to a communication channel, comprising the steps of:
  - (a) generating a parity check matrix comprising:

M tiers,

wherein  $M \ge 2$ ,

Dmin = 2 \* M for M = 1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

tc = M,

wherein to is the column weight, and

cycle -4 = 0;

- (b) linear block encoding the data in accordance with the matrix generated in step (a); and
  - (c) transmitting the data encoded in step (b) to the communication channel.
- 156 158. (Currently Amended) A computer program according to Claim 155 157, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 157 159. (Currently Amended) A computer program according to Claim 156 158, wherein the rank of said identity matrix of one of said tiers is mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 158 160. (Currently Amended) A computer program according to Claim 156 158, wherein said M tiers are arranged in increasing rank order.
- 159 161. (Currently Amended) A computer program according to Claim 158 160, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .

- 160 162. (Currently Amended) A computer program according to Claim 156 158, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .
- 161 163. (Currently Amended) A computer program according to Claim 160 162, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 162 164. (Currently Amended) A computer program according to Claim 156 158, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 163 165. (Currently Amended) A computer program according to Claim 156 158, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- $\frac{164}{166}$ . (Currently Amended) A computer program according to Claim  $\frac{158}{160}$ , wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, & \text{if } c \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, & \text{otherwise} \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 165 167. (Currently Amended) A computer program according to Claim 158 160, wherein M = 3, the number of rows =  $P_1 + P_2 + P_3$ , the number of columns =  $P_1 * P_2$ , dmin  $P_1 = 0$ , and to  $P_2 = 0$ .
  - 166 168. (Currently Amended) A computer program according to Claim 165

<u>167</u>, wherein a code rate =  $(P_1 + P_2 + P_1 - P_2 - P_3 + 2)/(P_1 + P_2)$ .

167 169. (Currently Amended) A method for generating a method parity check matrix for one of a low-density parity-check encoder and a low-density parity-check decoder, comprising the steps of:

providing M tiers of element,

wherein  $M \ge 2$ ,

selecting Dmin = 2 \* M for M = 1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

selecting tc = M,

wherein tc is the column weight, and

selecting cycle - 4 = 0.

- 168 170. (Currently Amended) A method according to Claim 167 169, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 169 171. (Currently Amended) A method according to Claim 168 170, further comprising the step of setting the rank of said identity matrix of one of said tiers to be mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 170 172. (Currently Amended) A method according to Claim 168 170, further comprising the step of arranging said M tiers in increasing rank order.
- 171 173. (Currently Amended) A method according to Claim 170 172, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
  - 172 174. (Currently Amended) A method according to Claim 168 170, wherein

said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .

- 173 175. (Currently Amended) A method according to Claim 172 174, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 174 176. (Currently Amended) A method according to Claim 168 170, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 175 177. (Currently Amended) A method according to Claim 168 170, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- $\frac{176}{178}$ . (Currently Amended) A method according to Claim  $\frac{170}{172}$ , wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, ifc \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, otherwise \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 177 179. (Currently Amended) A method according to Claim 170 172, wherein M = 3, the number of rows =  $P_1 + P_2 + P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and tc = 3.
- 178 180. (Currently Amended) A method according to Claim 177 179, wherein a code rate =  $(P_1 * P_2 + P_1 P_2 P_3 + 2)/(P_1 * P_2)$ .

179 181. (Currently Amended) A computer program embodied in a medium for generating a computer program parity check matrix for one of a low-density parity-check encoder and a low-density parity-check decoder, comprising the steps of:

providing M tiers of element,

wherein  $M \ge 2$ ,

selecting Dmin = 2 \* M for M=1..3 or

 $2*M \ge Dmin \ge 6$  for M > 3,

wherein Dmin is the minimum Hamming distance,

selecting tc = M,

wherein tc is the column weight, and selecting cycle - 4 = 0.

- 180 182. (Currently Amended) A computer program according to Claim 179 181, wherein each of said M tiers comprises an identity matrix having a corresponding rank  $P_i$ , wherein  $1 \le i \le M$ .
- 181 183. (Currently Amended) A computer program according to Claim 180 182, further comprising the step of setting the rank of said identity matrix of one of said tiers to be mutually prime with respect to the rank of said identity matrix of another one of said tiers.
- 182 184. (Currently Amended) A computer program according to Claim 180 182, further comprising the step of arranging said M tiers in increasing rank order.
- 183 185. (Currently Amended) A computer program according to Claim 182 184, wherein said matrix comprises C columns, wherein  $C \le P_1 * P_2$ .
  - 184 186. (Currently Amended) A computer program according to Claim 180

182, wherein said matrix comprises R rows, wherein  $R = \sum_{i=1}^{M} P_i$ .

- 185 187. (Currently Amended) A computer program according to Claim 184 186, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  independent rows.
- 186 188. (Currently Amended) A computer program according to Claim 180 182, wherein said matrix comprises  $\sum_{i=1}^{M} P_{i-1}(M-1)$  parity bits.
- 187 189. (Currently Amended) A computer program according to Claim 180 182, wherein said matrix comprises  $P_1 \times P_2 \sum_{i=1}^{M} P_i + (M-1)$  maximum user bits.
- $\frac{188}{190}$ . (Currently Amended) A computer program according to Claim  $\frac{182}{184}$ , wherein for each element  $A_{r,c}$

$$For \sum_{j=1}^{n-1} P_j + 1 \le r \le \sum_{j=1}^{n} P_j,$$

$$A_{r,c} = \begin{cases} 1, & \text{if } c \mod(P_n) = r - \sum_{j=1}^{n-1} P_j \\ 0, & \text{otherwise} \end{cases}$$

$$0 \le c \le C$$

$$C \le P_1 * P_2$$

- 189 191. (Currently Amended) A computer program according to Claim 182 184, wherein M = 3, the number of rows =  $P_1 + P_2 + P_3$ , the number of columns =  $P_1 * P_2$ , dmin Dmin = 6, and tc = 3.
- 190 192. (Currently Amended) A computer program according to Claim 189 191, wherein a code rate =  $(P_1 * P_2 + P_1 P_2 P_3 + 2)/(P_1 * P_2)$ .